

Performance Analysis of MAC Protocols for Wireless LANs in Rayleigh and Shadow Fading Channel

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Background

■ Why Wireless LAN ?

- Increase the wireless service.
- Easy to change the network configuration.
- Unlicensed Wireless Channel : ISM Band (900MHz, 2.4GHz, 5.7GHz)
- Increasing the wireless LANs market.

■ Wireless LAN Standards

- IEEE 802.11 : *CSMA/CA (Carrier Sense Multiple Access with Collision Avoidance) protocol*
- ETSI HIPERLAN Type I : EY-NPMA (Elimination Yield -Non-preemptive Priority) protocol



Previous Works

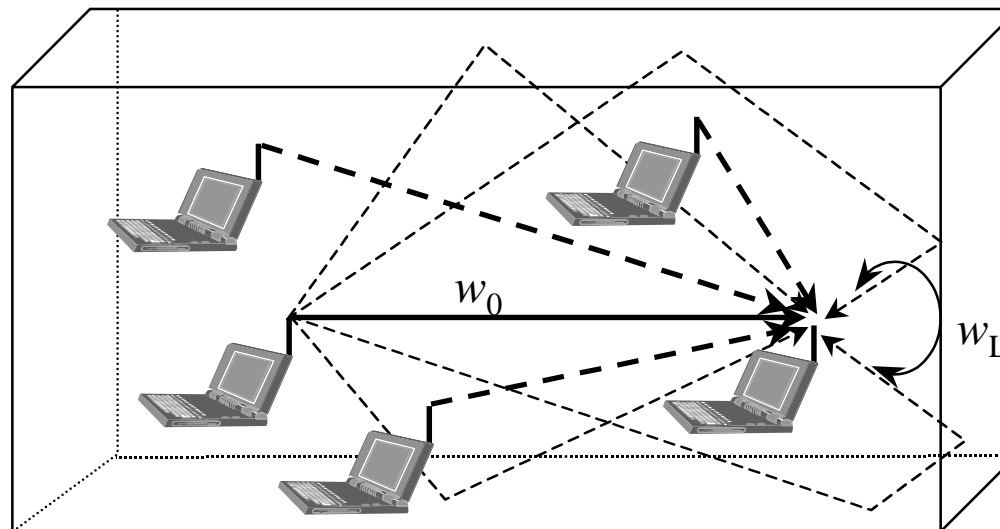
- K. C. Chen : Throughput and delay analysis of CSMA/CA (Personal, Indoor and Mobile Radio Communication: PIMRC '95)
 - Non-persistent CSMA => *No Backoff delay*
- S. M. Nor : Computer simulation for CSMA/CA (PIMRC '95)
 - *No analytical model*
- H. Chhaya and S. Gupta : CSMA/CA (PIMRC '95)
 - Calculate the lower bound of throughput=>*No Backoff delay*
- J. L. Sobrinho and A. Krishnakumar (GLOBECOM '96)
 - Consider the CSMA/CA for only real time service(voice)
 - *Delay stability analysis*
- Our Research
 - CSMA/CA : *1-persistent + p-persistent CSMA model*
 - *Wireless fading channel model*
 - *Throughput & Packet delay*



Radio Channel Model

■ Radio Channel (Linnartz model)

- Multipath fading : Rayleigh Distribution
- Shadow fading : Log-normal Distribution
- Near-Far Effect : Uniform Distribution



Radio Channel Model

■ Instantaneous power (w_0)

➤ *Unconditional probability density function of the instantaneous power*

$$f_{w_0}(w_0) = \int_0^\infty \int_0^\infty \underbrace{\frac{1}{w_L} \exp\left(-\frac{w_0}{w_L}\right)}_{\text{Multipath fading}} \underbrace{\frac{f(r_i)}{\sqrt{2\sigma_s w_L}} \exp\left\{-\frac{\ln^2(r_i^\xi w_L)}{2\sigma_s}\right\}}_{\text{Near-far effect}} \underbrace{\exp\left\{-\frac{\ln^2(r_i^\xi w_L)}{2\sigma_s}\right\}}_{\text{Shadowing}} dr_i dw_L$$

w_0 : the instantaneous power

w_L : the local mean power

r_i : a distance between a receiver and i -th transmitting terminal

ξ : path loss parameter (1.79 - 6.2, urban : 4)

σ_s : the standard deviation of Shadow fading

■ Capture Model

➤ Capture probability : $q(n|z)$

$$q(n|z) = \frac{2}{\sqrt{\pi}} \int_0^1 \int_{-\infty}^{\infty} r_1 \exp(-x_1^2) \left[\frac{1}{\sqrt{\pi}} \int_{-\infty}^{\infty} f(x_1, y_1) \exp(-y_1^2) dy_1 \right] dx_1 dr_1$$

$$\frac{2}{\sqrt{\pi}} \int_0^1 \int_{-\infty}^{\infty} r_2 \exp(-x_2^2) \left[\frac{1}{\sqrt{\pi}} \int_{-\infty}^{\infty} f(x_2, y_2) \exp(-y_2^2) dy_2 \right]^{n-1} dx_2 dr_2$$

where,

$$f(x_1, y_1) \triangleq \left[\sqrt{z} r_1^2 \exp\left\{ \frac{\sqrt{2}}{2} \sigma_s (x_1 - y_1) \right\} \right] \cdot \left[\frac{1}{z r_1^4} \exp\left\{ \frac{\sqrt{2}}{2} \sigma_s (x_1 - y_1) \right\} \right]$$

$$f(x_2, y_2) \triangleq \left[\sqrt{z} r_2^2 \exp\left\{ \frac{\sqrt{2}}{2} \sigma_s (x_2 - y_2) \right\} \right] \cdot \left[\frac{1}{z r_2^4} \exp\left\{ \frac{\sqrt{2}}{2} \sigma_s (x_2 - y_2) \right\} \right]$$

n : the number of stations,

z : capture ratio

r : a distance between terminals,

⇒ Use Hermite polynomial method for approximation

CSMA/CA Protocol

■ CSMA/CA

- Modified CSMA/CD (Carrier Sense Multiple Access/Collision Detection)
=> Can't detect collision in wireless channel

■ CSMA/CA packet transmission process

- Channel Idle : Immediately transmit
- Channel Busy : Wait until channel Idle => Backoff delay
=> If channel is Idle => Transmit
- Backoff delay algorithm
 - Increased exponentially according to the number of retransmission try.



Performance Analysis

■ Performance Parameters

- Throughput analysis
- Normalized packet delay analysis

■ CSMA/CA Protocols

- Basic CSMA/CA : Data- ...
- Stop-and-Wait CSMA/CA : Data-ACK-...
- 4-Way Hand Shake CSMA/CA : RTS-CTS-Data-ACK-...

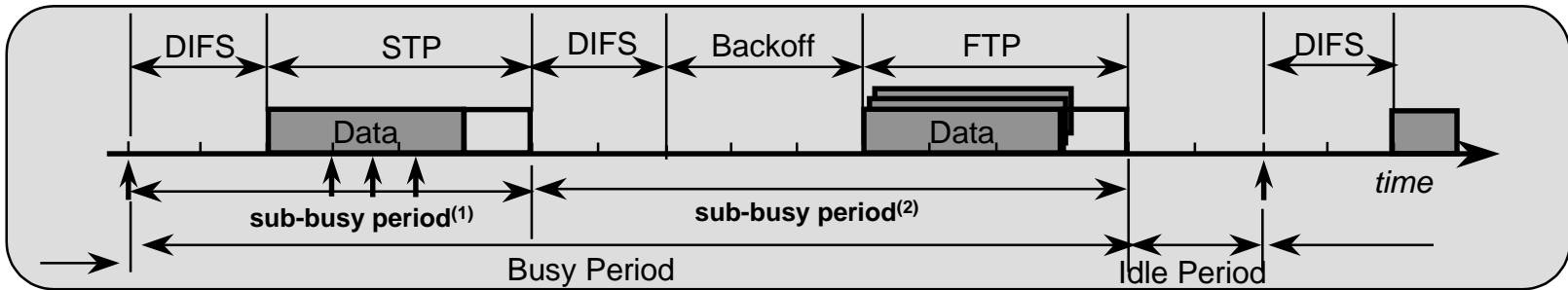
➤ ACK : Acknowledgement packet, RTS : Request To Send packet
CTS : Clear To Send packet



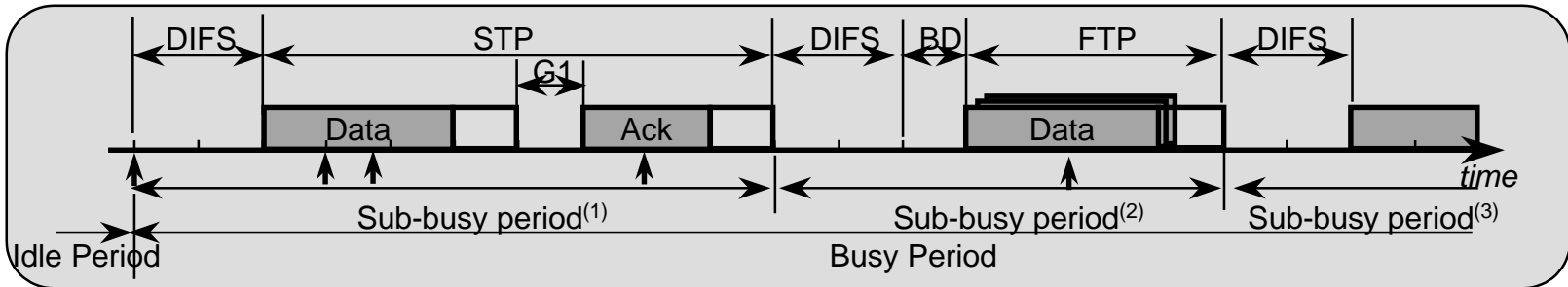
■ System Model

- CSMA/CA = 1-persistent CSMA + p -persistent CSMA
- M/G/1 Busy period analysis based on the Renewal Theory
- Finite population : The number of users (M)
- Slotted Channel: Slot size (a)
- All terminals are synchronized
- Packet arrival rate in a slot : g ($0 < g < 1$)
- Packet transmission probability : p ($0 < p \leq 1$)
- Channel model with Multipath fading, Shadow fading and Near-far effect

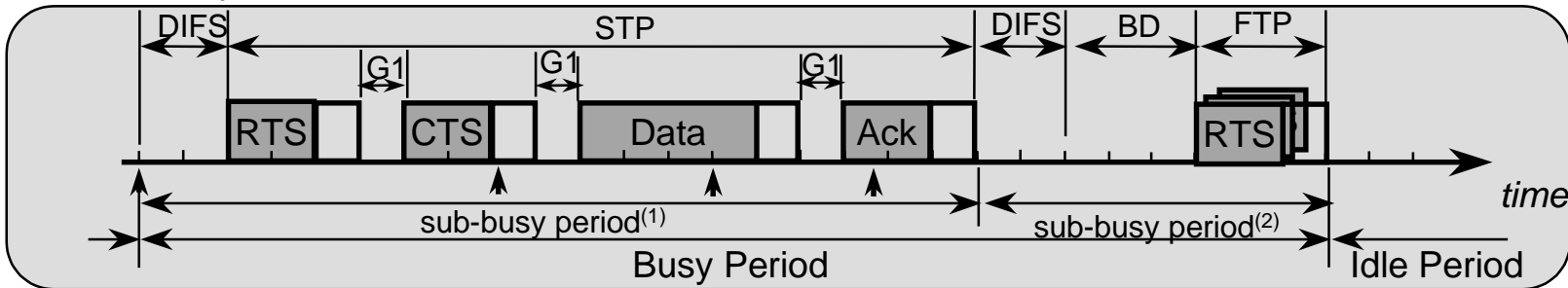
● Basic CSMA/CA Channel model



● Stop-and-Wait ARQ CSMA/CA Channel model



● 4-Way Handshake CSMA/CA Channel model



- DIFS : Distributed Inter-Frame Space
- STP : Successful Transmission Period
- RTS : Request To Send
- G1 : Short Inter-Frame Space
- FTP : Failed Transmission Period
- CTS : Clear To Send



■ CSMA/CA Throughput Analysis

- Basic concept

$$S = \frac{\overline{U}}{\overline{I} + \overline{B}}$$

- I : Idle period
- B: Busy period
- U : Useful transmission period



■ Throughput of Basic CSMA/CA

$$S = \frac{\overline{U}}{B + I}$$

$$\left[\sum_{i=1}^M \left\{ \frac{\binom{M}{i} [1 - (1-g)]^i (1-g)^{M-i}}{1 - (1-g)^M} \right\} \boxed{q(i, z)} + \frac{1}{1 - (1-g)^{((1/a)+1)M}} \sum_{n=1}^M \left[\left\{ \sum_{i=1}^n \binom{n}{i} p^i (1-p)^{n-i} \right. \right. \right.$$

$$\left. \left. \cdot \sum_{l=1}^n \binom{M-n}{l} g^l (1-g)^{M-n-l} + (M-n)(1-p)^{(k+1)n} g (1-g)^{(k+1)(M-n)-1} \boxed{(i+l)q(i+l, z)} \right\} \right.$$

$$\left. \cdot \frac{(1-p)^n (1-g)^{M-n}}{1 - (1-p)^n (1-g)^{M-n}} + \sum_{i=1}^n \binom{n}{i} p^i (1-p)^{n-i} \boxed{iq(i, z)} \right]$$

$$\cdot \left\{ \frac{\binom{M}{n} [1 - (1-g)^{(TP_s/a)]^n (1-g)^{(TP_s/a)(M-n)}}{1 - (1-g)^{(TP_s/a)M}} \right\}$$

$$= \left[\frac{f[1 - (1-g)^M] + 1 + a + \frac{a}{1 - (1-g)^{((1/a)+1)M}} \left(- (f+1+a)[1 - (1-g)^{((1/a)+1)M}] + \right. \right.$$

$$\left. \left. a \sum_{k=1}^{\infty} \{ (1-p)^k - (1-g)^{((1/a)+1)} [(1-p)^k - (1-g)^k] \}^M - a(1-g)^{((1/a)+a)M} \sum_{k=1}^{\infty} (1-g)^{kM} \right) \right.$$

$$\left. + \frac{a}{1 - (1-g)^M} \right]$$

$TP_s = 1+a+f$: Average period of successful transmission $\boxed{}$ capture probability

■ Throughput of SW ARQ & 4-WH CSMA/CA

$$S = \frac{\bar{U}}{\bar{B} + \bar{I}}$$

where

$$\bar{I} = \frac{a}{1 - (1 - g)^M}$$

$$\bar{B} = f \left[1 - (1 - g)^M \right]$$

$$+ \left\{ TP_{Succ} + [1 - (1 - g)^{(TP_{Succ}/a)}] B(TP_{Succ} / a) \right\} \bar{u}$$

$$+ \left\{ TP_{Fail} + [1 - (1 - g)^{(TP_{Fail}/a)}] B(TP_{Fail} / a) \right\} [1 - \bar{u}]$$

$$\bar{U} = \left\{ 1 + [1 - (1 - g)^{(TP_{Succ}/a)}] U(TP_{Succ} / a) \right\} \bar{u}$$

$$+ \left\{ [1 - (1 - g)^{(TP_{Fail}/a)}] U(TP_{Fail} / a) \right\} [1 - \bar{u}]$$

$$\text{SW ARQ CSMA/CA} : \begin{cases} TP_{Succ} = 1 + 2a + \beta + \delta + f & (\text{average period of successful transmission}) \\ TP_{Fail} = 1 + a + f & (\text{average period of failed transmission}) \end{cases}$$

$$4\text{-WH CSMA/CA} : \begin{cases} TP_{Succ} = 1 + \gamma + \theta + 3\beta + \delta + 4a + f \\ TP_{Fail} = \gamma + a + f \end{cases}$$

$$\bar{u} = [U^{(j)}]$$



■ CSMA/CA Delay Analysis

- Use channel throughput calculations
- Average number of retransmission for a packet

$$\left(\frac{G}{S} - 1 \right)$$

- Average delay for a packet from arrival time to successful transmission starting time (\bar{R})

- probability with arriving in Idle period : $\frac{\bar{I}}{\bar{B} + \bar{I}}$
- probability with arriving in Delay period : $\frac{\bar{D}}{\bar{B} + \bar{I}}$
- probability with arriving in Transmitting period : $\frac{\bar{T}}{\bar{B} + \bar{I}}$

- For Basic CSMA/CA

$$\bar{R} = \frac{\bar{I}}{\bar{B} + \bar{I}} f + \frac{\bar{D}}{\bar{B} + \bar{I}} f + \frac{\bar{T}}{\bar{B} + \bar{I}} \left[\frac{(1 + a + f + E[D^{(2)}])^2}{2(1 + a + f + E[D^{(2)}])} \right]$$

- \bar{Y} : Average Random delay by Backoff algorithm

■ Normalized packet delay for Basic CSMA/CA

$$L = \left(\frac{G}{S} - 1 \right) [1 + a + \bar{Y} + \bar{R}] + [\bar{R} + 1 + a]$$

■ Normalized packet delay for Stop-and-Wait CSMA/CA

$$L = \left(\frac{G}{S} - 1 \right) [1 + a + \bar{Y} + \bar{R}] + [\bar{R} + 1 + 2a + \beta + \delta]$$

■ Normalized packet delay for 4-Way Handshake CSMA/CA

$$L = \left(\frac{G}{S} - 1 \right) [\gamma + a + \bar{Y} + \bar{R}] + [\bar{R} + 1 + \gamma + \theta + \delta + 3\beta + 4a]$$

L = Normalized packet delay a = propagation delay,

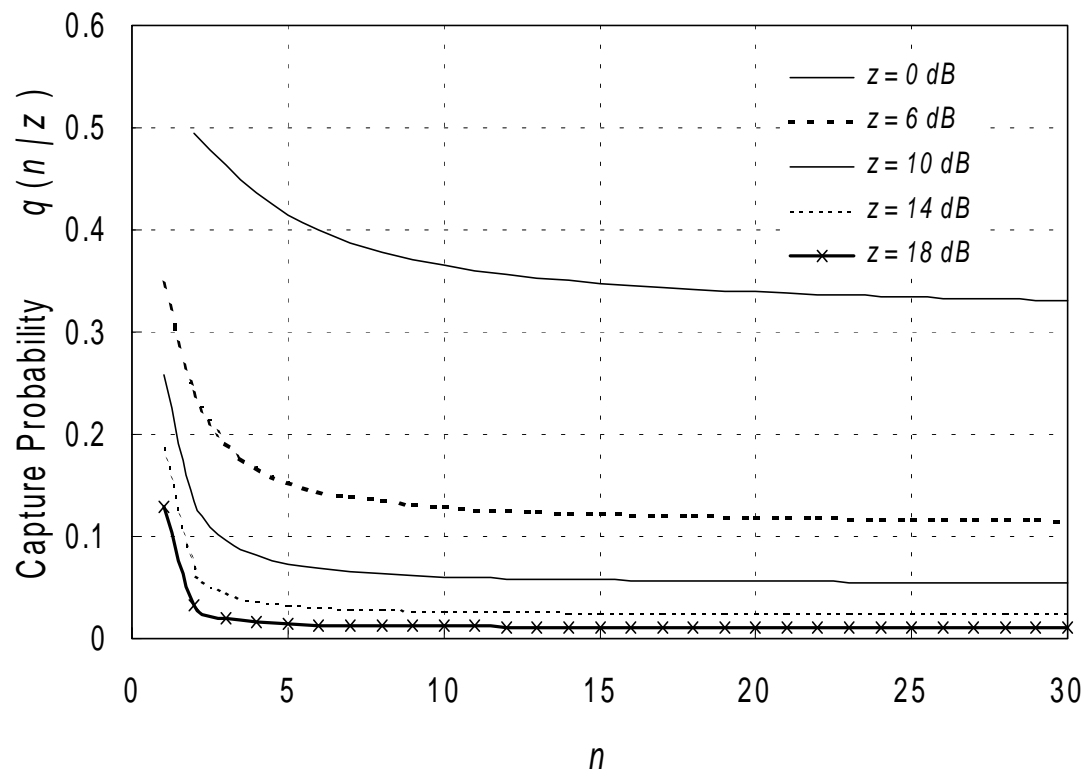
G = Offered load β = SIFS,

S = Throughput δ = ACK packet, γ = RTS packet, θ = CTS packet



Numerical Results and Simulations

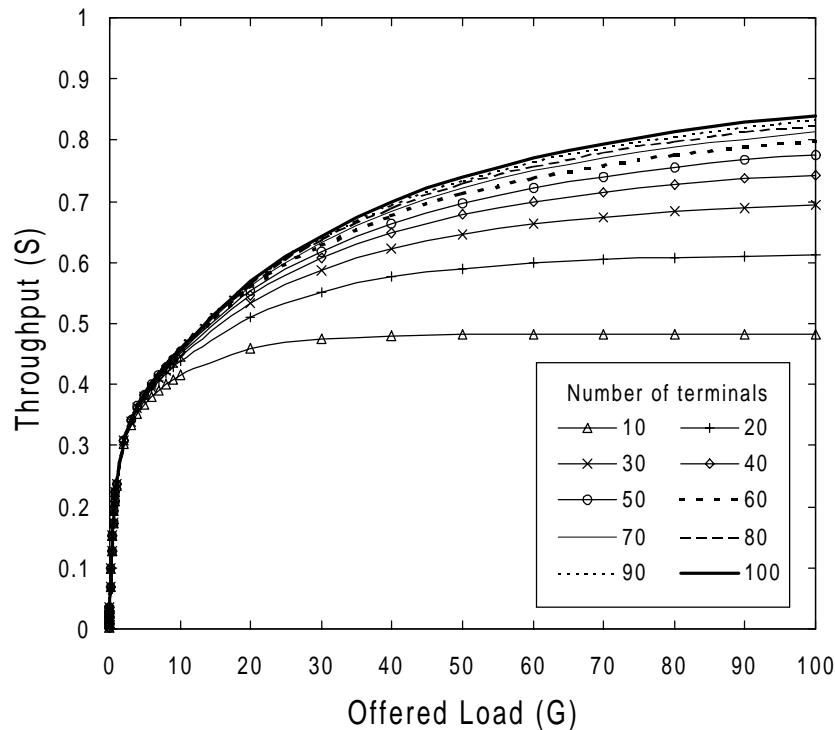
- Capture probability for the number of colliding pack



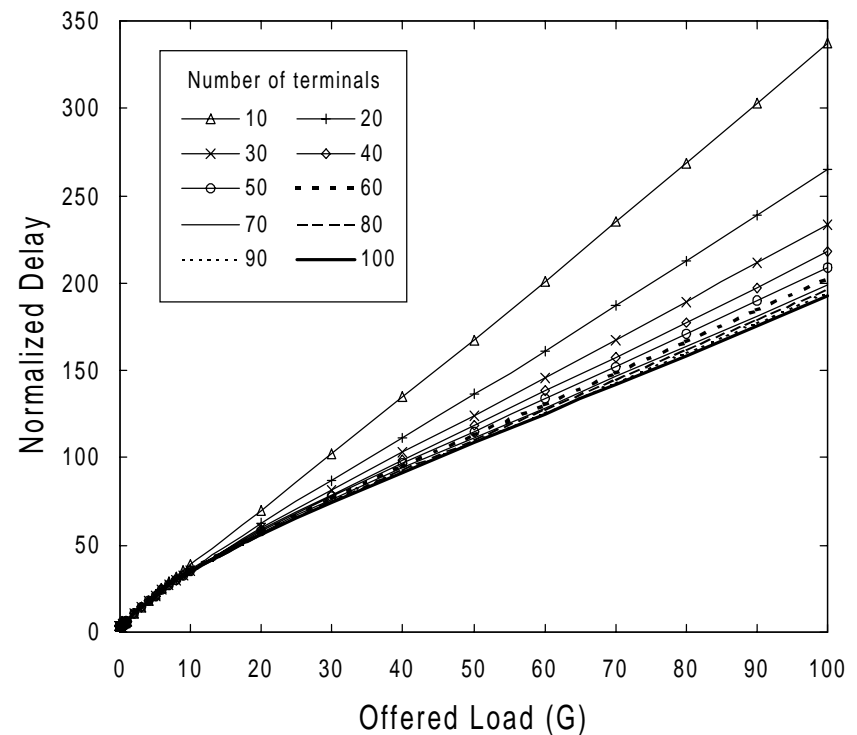
- standard deviation for Shadowing (σ_s) = 6dB
- attenuation factor for the distance (ξ) = 4

■ Throughput and Packet delay vs. offered load of Basic CSMA/CA for varying of the number of terminals

(a) Throughput



(b) Packet delay



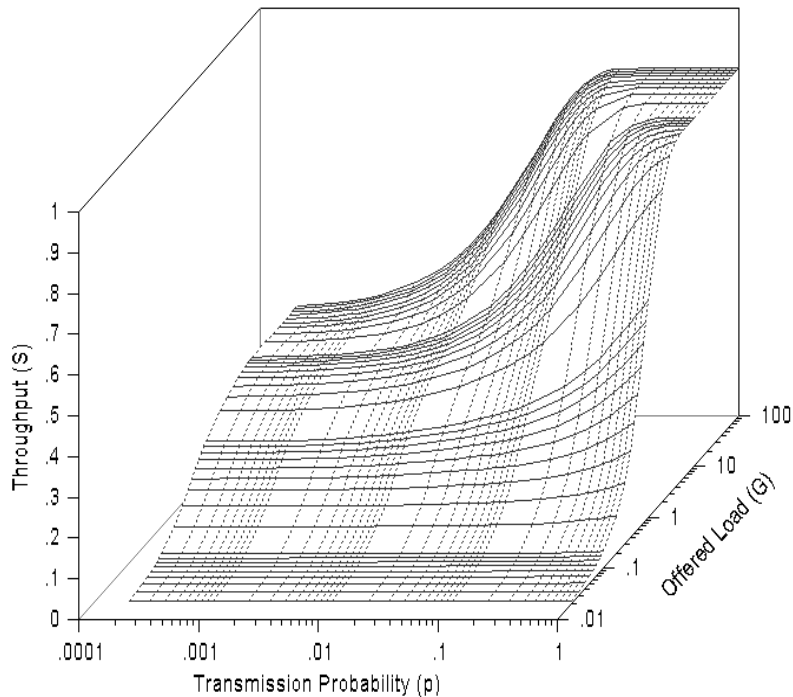
- propagation delay (a) = 0.01
- packet transmission probability (p) = 0.03,
- standard deviation for shadowing (σ_s) = 6 dB
- attenuation factor for distance (ξ) = 4

- DIFS length (f) = 0.06
- random delay ($E(Y)$) = 0.06
- capture ratio (z) = 4

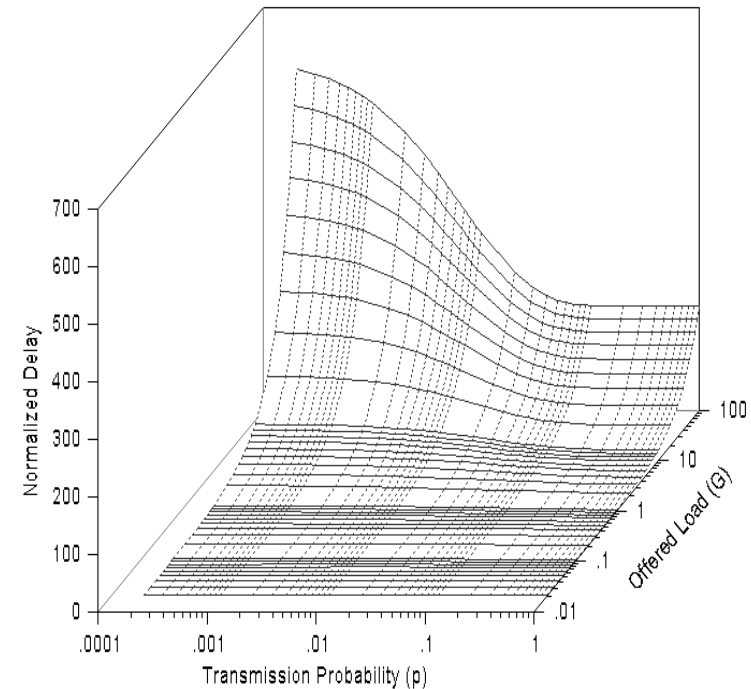


■ Throughput and Packet delay vs. transmission probability p of SW ARQ CSMA/CA for varying of the offered load

(a) Throughput



(b) Packet delay

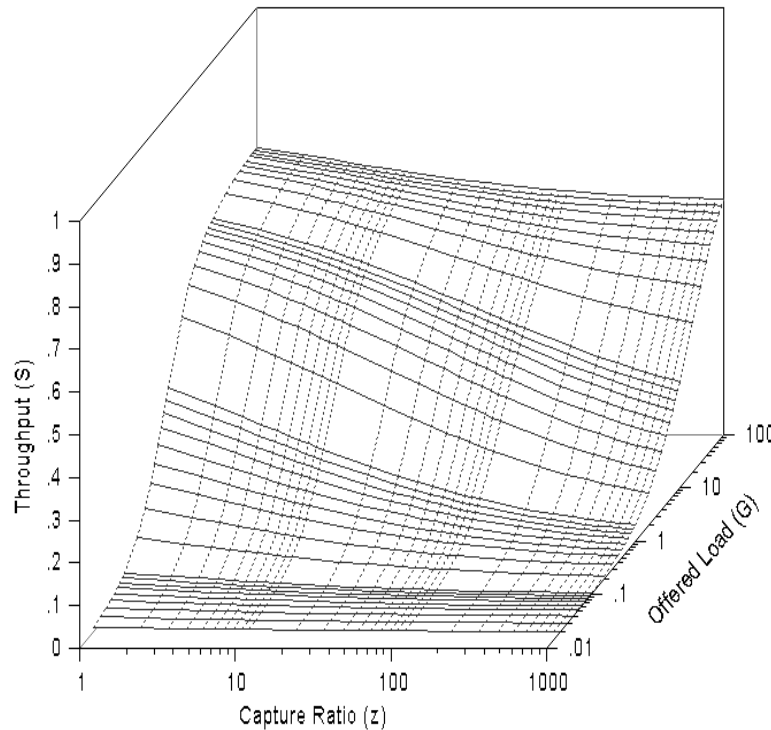


- propagation delay (a) = 0.01
- standard deviation for shadowing (σ_s) = 6 dB
- attenuation factor for distance (ξ) = 4
- number of terminals = 50
- short Inter-Frame Space = 0.03

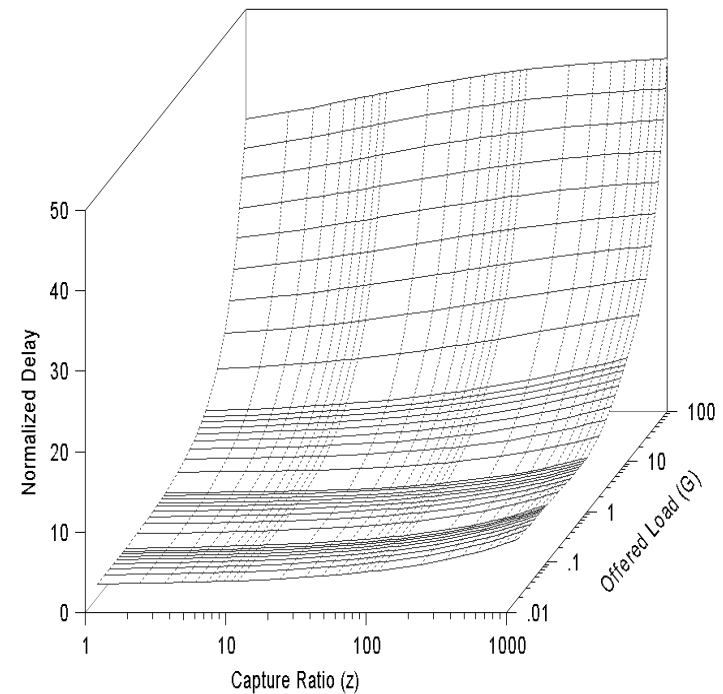
- DIFS length (f) = 0.06
- random delay ($E(Y)$) = 0.06
- capture ratio (z) = 4
- length of Ack packet = 0.06

■ Throughput and Packet delay vs. capture ratio z of 4-WH CSMA/CA for varying of the offered load

(a) Throughput



(b) Packet delay

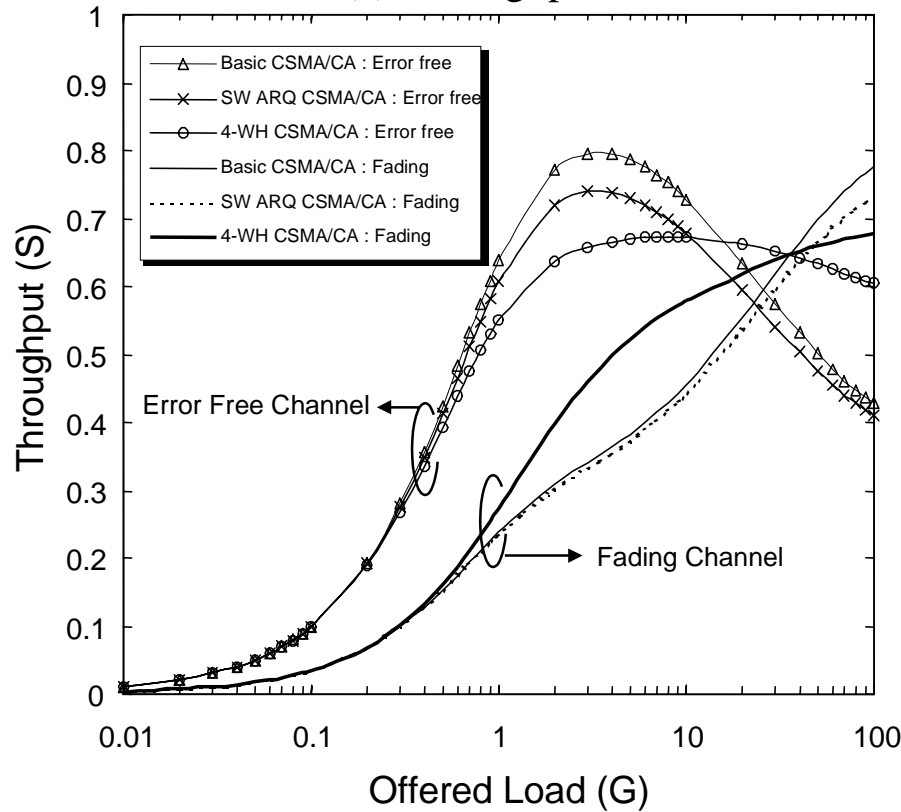


- propagation delay (a) = 0.01
- packet transmission probability (p) = 0.03
- standard deviation for shadowing (σ_s) = 6 dB
- attenuation factor for distance (ξ) = 4
- number of terminals = 20
- random delay ($E(Y)$) = 0.06

- length of DIFS (f) = 0.06
- length of Ack packet = 0.06
- length of SIFS = 0.03
- length of RTS = 0.06
- length of CTS = 0.03

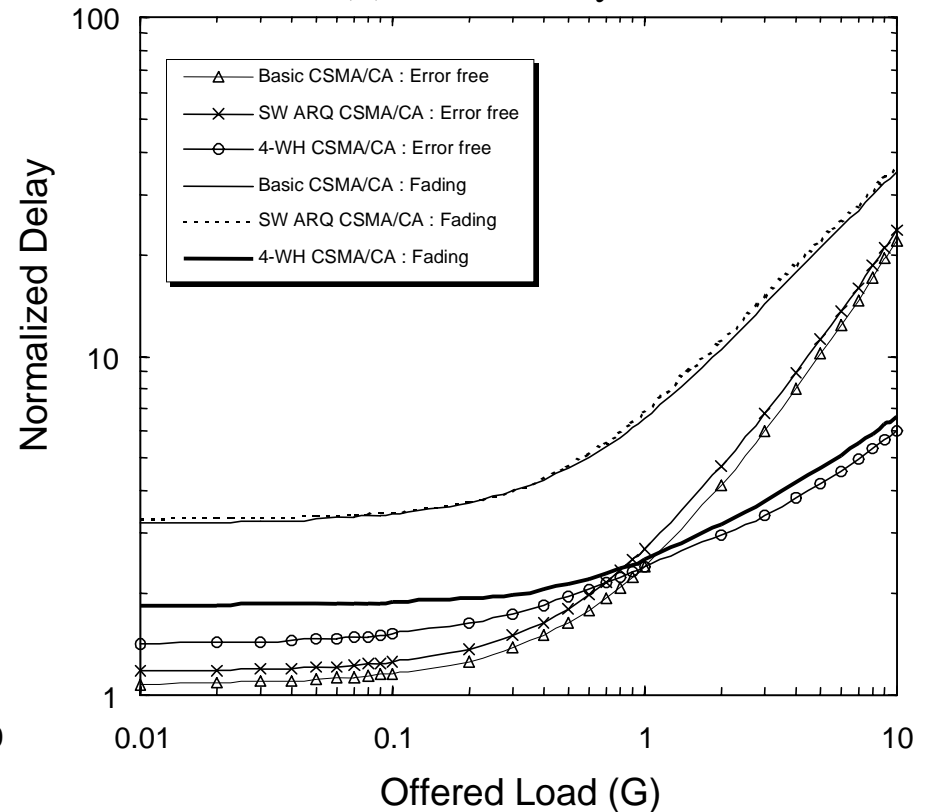
■ Performance comparison of three types of CSMA/CA protocols for error free channel model and fading channel model

(a) Throughput



- propagation delay (a) = 0.01
- packet transmission probability (p) = 0.03
- standard deviation for shadowing (σ_s) = 6 dB
- attenuation factor for distance (ξ) = 4
- capture ratio = 4
- number of terminals = 50

(b) Packet delay



- length of DIFS (f) = 0.06
- length of Ack packet = 0.06
- length of SIFS = 0.03
- length of RTS = 0.06
- length of CTS = 0.03
- random delay ($E(Y)$) = 0.06

Conclusions

- Propose a new analytical approach for CSMA/CA protocols in *Multipath fading, Shadow fading and Near-far Effect* environments
- Throughput and packet delay analysis
- Performance comparison of three type CSMA/CA protocols with and without *radio channel model*
- Performance prediction for various parameters
- Further Study
 - Integration of PCF(Point Coordination Function)
 - HIPERLAN : EY-NPMA protocol

■ Probability of successful transmission for SW ARQ & 4-WH CSMA/CA

$$E[U^{(j)}] = \begin{cases} u(1) & ; j = 1 \\ u(\text{STP}/a) & ; \text{if } (j-1)\text{th transmission is success, } j = 2, 3, \dots \\ u(\text{FTP}/a) & ; \text{if } (j-1)\text{th transmission is unsuccess, } j = 2, 3, \dots \end{cases}$$

where

$$u(1) = \sum_{i=1}^M \left\{ \frac{\binom{M}{i} [1 - (1-g)]^i (1-g)^{M-i}}{1 - (1-g)^M} \right\} q(i-1, z)$$

$$u(\text{STP}/a) = \sum_{n=0}^M \left[\left\{ \sum_{i=1}^n \binom{n}{i} p^i (1-p)^{n-i} \sum_{l=1}^n \binom{M-n}{l} g^l (1-g)^{M-n-l} \right\} (i+l)q(i+l, z) \right. \\ \left. \frac{(1-p)^n (1-g)^{M-n}}{1 - (1-p)^n (1-g)^{M-n}} + \sum_{i=1}^n \binom{n}{i} p^i (1-p)^{n-i} \right\} i q(i-1) \right] \\ \cdot \left\{ \frac{\binom{M}{n} [1 - (1-g)^{(\text{STP}/a)}]^n (1-g)^{(\text{STP}/a)(M-n)}}{1 - (1-g)^{(\text{STP}/a)M}} \right\}$$

$$u(\text{FTP}/a) = \sum_{n=0}^M \left[\left\{ \sum_{i=1}^n \binom{n}{i} p^i (1-p)^{n-i} \sum_{l=1}^n \binom{M-n}{l} g^l (1-g)^{M-n-l} \right\} (i+l)q(i+l, z) \right. \\ \left. \frac{(1-p)^n (1-g)^{M-n}}{1 - (1-p)^n (1-g)^{M-n}} + \sum_{i=1}^n \binom{n}{i} p^i (1-p)^{n-i} \right\} i q(i-1) \right] \\ \cdot \left\{ \frac{\binom{M}{n} [1 - (1-g)^{(\text{FTP}/a)}]^n (1-g)^{(\text{FTP}/a)(M-n)}}{1 - (1-g)^{(\text{FTP}/a)M}} \right\}$$

STP : Successful Transmission period, FTP : Failed Transmission period

$\{\}$ capture probability

